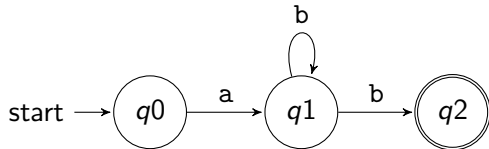
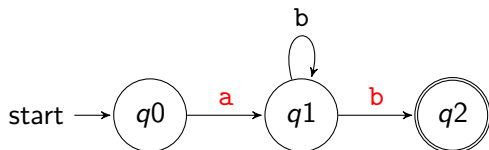


Finite automata



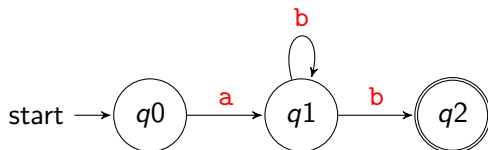
Finite automata



accepting runs $q_0 \xrightarrow{a} q_1 \xrightarrow{b} q_2 \in F$

ab

Finite automata

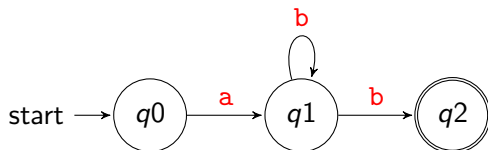


accepting runs $q_0 \xrightarrow{a} q_1 \xrightarrow{b} q_2 \in F$

$q_0 \xrightarrow{a} q_1 \xrightarrow{b} q_1 \xrightarrow{b} q_2 \in F$

ab, abb

Finite automata



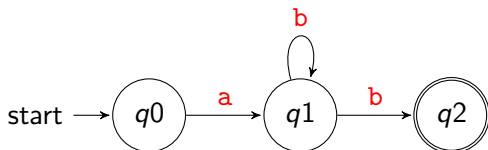
accepting runs $q_0 \xrightarrow{a} q_1 \xrightarrow{b} q_2 \in F$

$q_0 \xrightarrow{a} q_1 \xrightarrow{b} q_1 \xrightarrow{b} q_2 \in F$

\vdots

regular expression ab^*b denoting the regular language $\{ab, abb, \dots\}$

Finite automata and regular grammars: example



accepting runs $q_0 \xrightarrow{a} q_1 \xrightarrow{b} q_2 \in F$

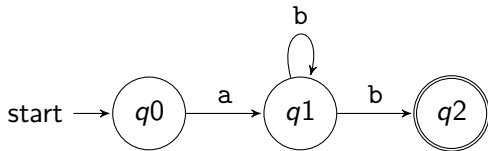
$q_0 \xrightarrow{a} q_1 \xrightarrow{b} q_1 \xrightarrow{b} q_2 \in F$

\vdots

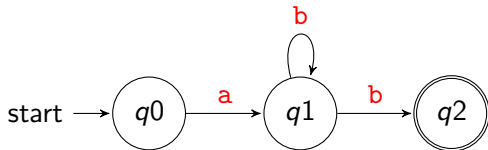
regular expression ab^*b denoting the regular language $\{ab, abb, \dots\}$

```
q0 --> [a], q1.      % tran(q0,a,q1).
q1 --> [b], q2.      % tran(q1,b,q2).
q1 --> [b], q1.      % tran(q1,b,q1).
q2 --> [].           % final(q2).
```

Strings accepted by a finite automaton



Strings accepted by a finite automaton

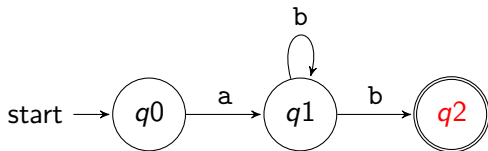


$\text{tran}(q_0, a, q_1).$

$\text{tran}(q_1, b, q_2).$

$\text{tran}(q_1, b, q_1).$

Strings accepted by a finite automaton



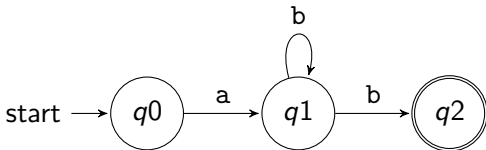
`tran(q0,a,q1).`

`tran(q1,b,q2).`

`tran(q1,b,q1).`

`final(q2).`

Strings accepted by a finite automaton



`tran(q0,a,q1).`

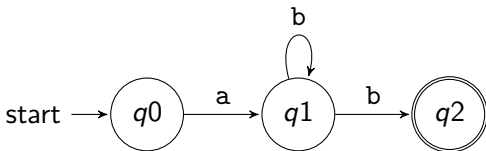
`tran(q1,b,q2).`

`tran(q1,b,q1).`

`final(q2).`

`accept(String) :- steps(q0,String,Q), final(Q).`

Strings accepted by a finite automaton



`tran(q0,a,q1).`

`tran(q1,b,q2).`

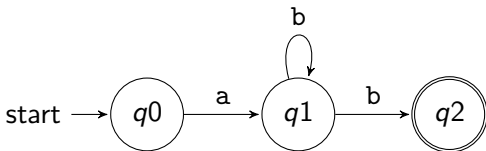
`tran(q1,b,q1).`

`final(q2).`

`accept(String) :- steps(q0,String,Q), final(Q).`

`steps(Q,[H|T],N) :- tran(Q,H,Qn), steps(Qn,T,N).`

Strings accepted by a finite automaton



`tran(q0,a,q1).`

`tran(q1,b,q2).`

`tran(q1,b,q1).`

`final(q2).`

`accept(String) :- steps(q0,String,Q), final(Q).`

`steps(Q,[],Q).`

`steps(Q,[H|T],N) :- tran(Q,H,Qn), steps(Qn,T,N).`

Regular expressions over a finite alphabet A

- ▶ \emptyset denoting the empty set
- ▶ ϵ denoting the set $\{\epsilon\}$ containing only the empty string, which has no characters at all
- ▶ a in A denoting the set $\{a\}$

Regular expressions over a finite alphabet A

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and for any regular expressions E and F ,

- ▶ concatenation (EF) denoting the set of strings made by concatenating a string from E with a string from F
- ▶ alternation ($E + F$) denoting the union of sets described by E and by F (also written $E|F$)
- ▶ Kleene star (E^*) denoting the set of all strings made by concatenating any finite number (including zero) of strings from E .