



## B<sup>+</sup>-Trees

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## B<sup>+</sup>-Trees

most common variant of B-tree index is B<sup>+</sup>-tree  
data pointers are stored *only at the leaf nodes*  
structure of leaf nodes differs from internal nodes  
leaf nodes have entry for *every* value of the search field  
along with a pointer to the record (or the block that  
contains the record)

for non-key search field, pointer points to a block  
containing pointers to the data file records, creating  
extra level of indirection

leaf nodes generally linked together to provide ordered  
access on the search field to the records

B<sup>+</sup>-Trees

2

## B<sup>+</sup>-Trees

*Structure of internal nodes of a B<sup>+</sup>-tree:*

each internal node is of the form:

$\langle A_1, K_1, A_2, K_2, \dots, A_{k-1}, K_{k-1}, A_k \rangle$

where  $k \leq m$  and each  $A_i$  is a tree pointer

within each internal node,  $K_1 < K_2 < \dots < K_{k-1}$

for all search field values  $X$  in the sub-tree pointed at by  
 $A_i$ , we have

$K_{i-1} < X < K_i$  for  $1 < i < k$ ,

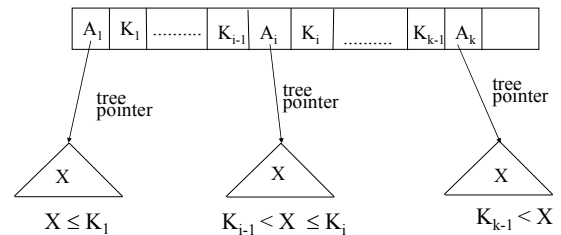
$X \leq K_1$  for  $i=1$ , and

$K_{k-1} < X$  for  $i=k$

B<sup>+</sup>-Trees

3

## Structure of Internal nodes of B<sup>+</sup>-Trees



Internal node of a B<sup>+</sup>-tree

B<sup>+</sup>-Trees

4

## Structure of Internal nodes of B<sup>+</sup>-Trees

each internal node has at most  $m$  tree pointers

each internal node, except the root, has at least  $\text{CEIL}[m/2]$  tree pointers; the root node has at least 2 tree pointers if it is an internal node

an internal node with  $k$  pointers,  $k \leq m$ , has  $k-1$  search field (key) values.

B<sup>+</sup>-Trees

5

## Structure of Leaf nodes of B<sup>+</sup>-Trees

Each leaf node is of the form

$\langle \langle K_1, R_1 \rangle, \langle K_2, R_2 \rangle, \dots, \langle K_{k-1}, R_{k-1} \rangle, R_{\text{next}} \rangle$

where  $k \leq m$ , each  $R_i$  is a data pointer, and  $R_{\text{next}}$  points to the next leaf node of the B<sup>+</sup>-tree

within each leaf node,  $K_1 < K_2 < \dots < K_{k-1}$ ,  $k \leq m$

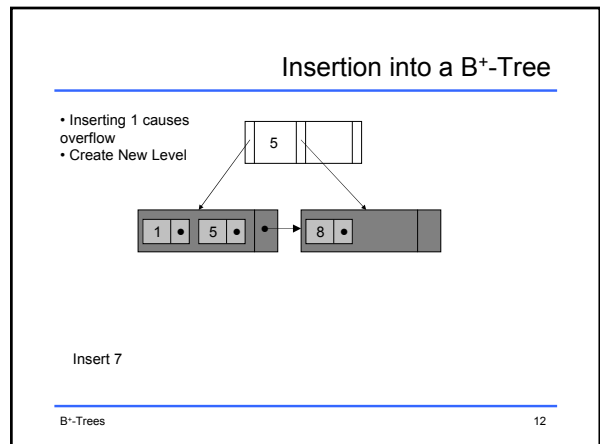
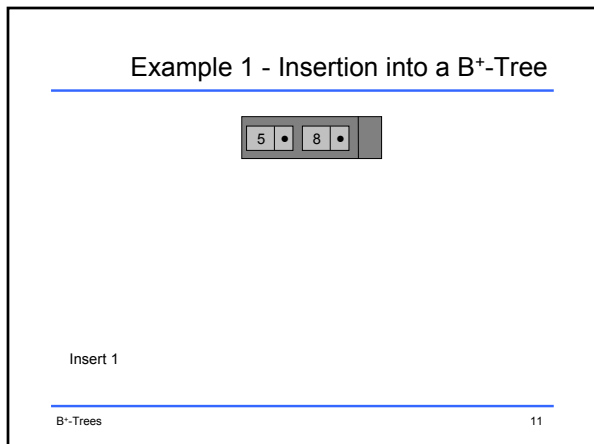
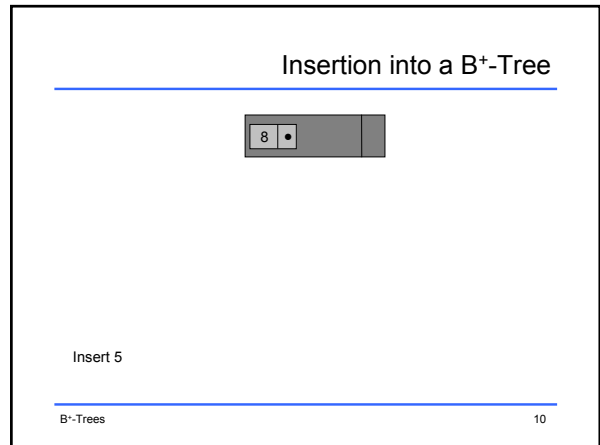
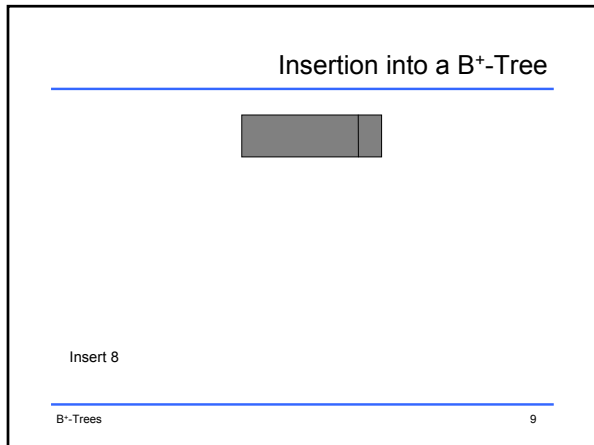
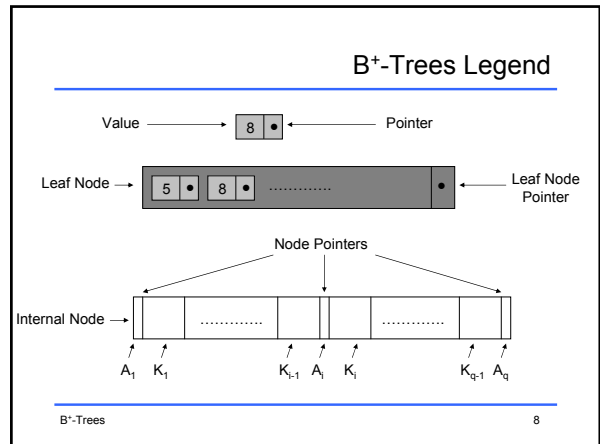
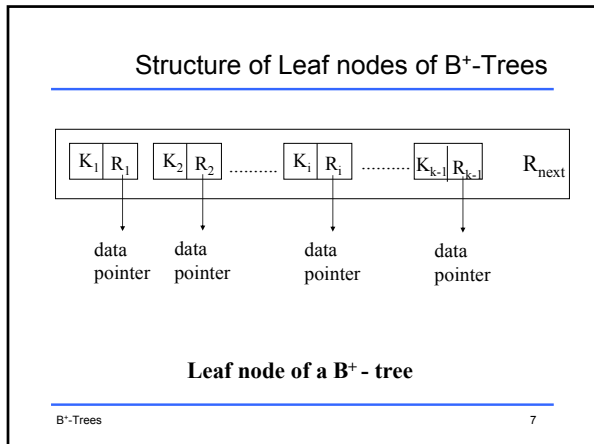
Each  $R_i$  is a *data pointer* that points to the records whose search field value is  $K_i$  or to a file block containing the record (or a block of record pointers that point to records whose search field value is  $K_i$  if the search field is not a key)

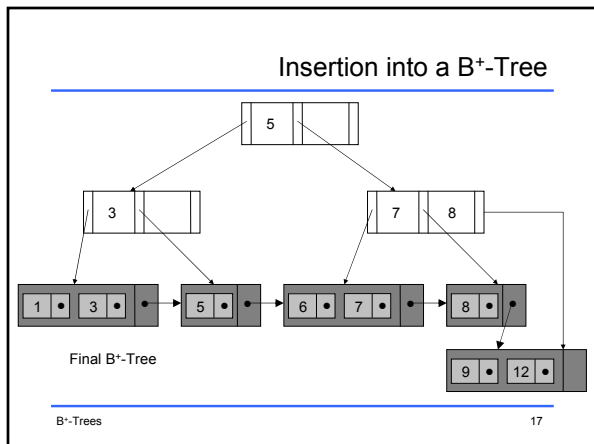
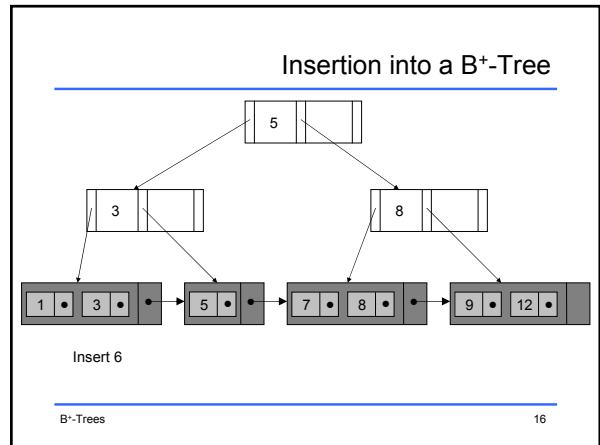
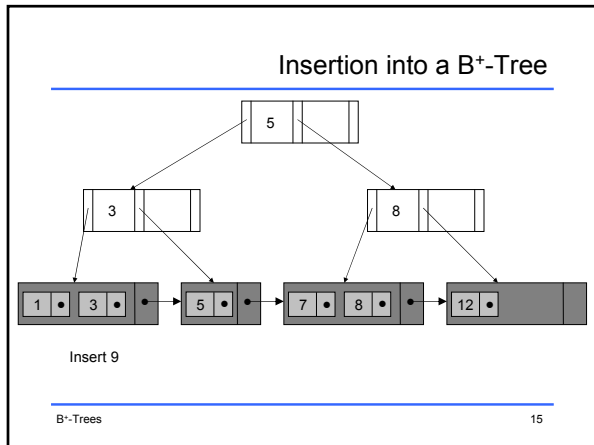
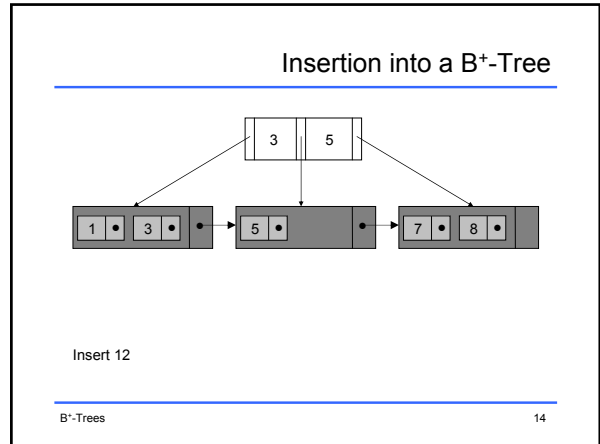
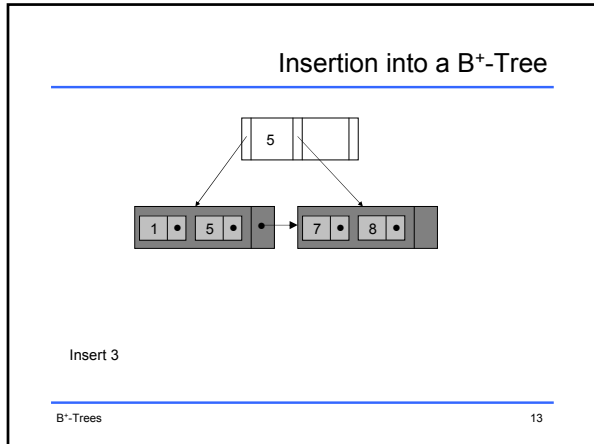
Each leaf node has at least  $\text{CEIL}[m/2] - 1$  values

All leaf nodes are at the same level

B<sup>+</sup>-Trees

6





- ### B<sup>+</sup>-Tree Variations
- Many variations
    - B-trees
    - B<sup>+</sup>-trees
    - B<sup>\*</sup>-trees
  - Common modifications
    - Change the fillfactor from 0.5 to 1.0
    - Allow a node to become empty before merging
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- B<sup>+</sup>-Trees 18

## Review

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B<sup>+</sup>-Tree is the most common variant of the B-Tree

It has two types of nodes

- Internal nodes
- Leaf nodes

Is more efficient than a B-Tree (see handout)